

# Load balancing in cloud using prioritization based on Quality of Services (QoS) demand

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**Abstract**— Cloud computing is the emerging technology that delivers the services on demand using the underline internet based software's. It is used to store a large amount of data and provide the accessibility to that data with minimum response time. Minimum bandwidth is required to assess the stored data from anywhere. In cloud there is a shared repository of resources that the cloud service providers try to utilize efficiently. Cloud vendors lease these services to the cloud customers as per their needs or demands. Customers also pay to the vendors as per their usage of cloud services. It is beneficial for both the service provider and customer as service provider earn money leasing their resources and customers also earn using the resources which otherwise they need to purchase. Cloud computing generally uses the concepts of virtual machines (VM) and their scheduling to feed maximum customers. This paper briefs about how to improve the quality of the services provided to the customer. Here we assign priorities to the customers on the basis of certain factors like number of VMs required, request type based on pricing and association with company etc. On the basis of these factors we assign priorities to the customers and the customers are feed on the basis of this priority. Here we also included the concepts of feeding the users request from the same geographical areas with minimum distance. This will improve the overall performance and throughput of the cloud with maximum customer satisfaction.

**Keywords**— Cloud Computing, Load balancing, Priority , Laod balancer

## I. INTRODUCTION

Cloud computing is an exceptionally current theme and the term has picked up a great deal of consideration as of late. It can be characterized as on request pay according to utilize display in which shared assets, data, programming and different gadgets are given by the customers' prerequisite when required. Human reliance on cloud is clear from the way that today's most well-known social organizing, email, archive sharing and web based gaming destinations are facilitated on cloud. Google, Microsoft, IBM, Amazon, Yahoo and Apple among others are exceptionally dynamic in this field [1].

Load balancing is the significant zone of worry in cloud computing. It is a component that disperses the workload equitably over every one of the hubs in the entire cloud to maintain a strategic distance from a circumstance where a few hubs are vigorously stacked while others are sit still or doing little work. Load balancing is the way toward reassigning the aggregate burdens to the individual hubs of the aggregate framework to make the best reaction time and furthermore great usage of the assets. Cloud computing is an administration giving office over the web in which the heap adjusting is the one of the testing undertaking. Different

strategies are to be utilized to improve a framework by distributing the loads to the hubs in an adjusting way however because of system blockage, transmission capacity use and so forth, there were issues are happened [1][2]. These issues were tackled by a portion of the current algorithms. A load balancing algorithm which is alterable in nature does not require any past data conduct of the framework, that is, it relies on upon the present conduct of the framework. There are different objectives that identified with the load balancing, for example, to enhance the execution significantly, to keep up the framework steadiness and so on. Contingent upon the present condition of the framework, load balancing calculations can be arranged into two sorts they are static and dynamic algorithms.

## II. LOAD BALANCING

It disseminates the aggregate load to the individual machines of the framework in such a way that each and every hub successfully uses the assets and limits the reaction time.

It tackles situations where a portion of the hubs are over stacked or under stacked. A dynamic calculation does not consider the past state & relies upon the present conduct of the framework [3]. The key qualities utilized are: execution of

framework, correlation of load, security of various framework, collaboration between the hubs, estimation of load, choosing of hubs, nature of work to be exchanged and numerous different ones .Load balancing can be categories into basic two types:

#### A. Static load balancing:

Here the cloud supplier uses homogeneous assets. Here progressed provisioning is done in which the suppliers set up the suitable assets before beginning of administrations as per the agreements done. Non versatile assets are used due to the static condition. The cloud requires prior information of hubs limit, preparing power, memory, execution and insights of client necessities which are not variable [4].

#### B. Dynamic load balancing:

Heterogeneous assets are supplied by cloud here. Dynamic provisioning is done with adaptable assets in which the specialist co-op apportionments more assets as they are expected to the client and expelled them when they are not utilized. Cloud does not require any previous learning [5] and the prerequisites of the clients are versatile. Dynamic condition is hard to be reproduced yet is very customizable with distributed computing condition.

### III. METHODOLOGY

It uses the priority which is assigned to both user tasks and datacenters depending on parameters stated. 5.1 Prioritization on user: After submitting the tasks, it will be processed by LB which calculates the priority. The user task will provide parameters like cost based on QoS request (CBQR) provided 60% priority which satisfies SLA conditions which aims to increase the overall efficiency [3][8]. More the amount of CBQR, high the priority it will be assigned. Second parameter will be the count of VMs required which is assigned 40% priority. Less priority will be assigned to the task which will require more number of VM. Now the load balancer will calculate the total priority of the task using table 1 and 2.

5.2 Prioritization on datacenters: Again the prioritization to datacenters will be assigned by the load balancers on the basis of one parameter. More number of VMs a datacenter possess, higher priority it will be assigned. The table of the datacenters priority (table 3) will be available to load balancer [4]. On the basis of this the load balancer will prepare table of datacenters on the basis of their geographical location (GL) and the count of VMs they have and also the number of available VMs in a datacenter.

5.3 Load balancer: It is the middleware which is responsible for assigning the priority to user tasks and datacenters. It is

the heart of the algorithm & contains table on CBQR which describes the range of priority.

CBQR		
Quality	Cost Range	Priority
Excellent	$CBQR > 1000000$	60%
Good	$50000 < CBQR \leq 1000000$	30%
Average	$10000 < CBQR \leq 50000$	10%

Table 1: Priority assignment based on CBQR

It also allocates priority to a parameter that is VMs count required. More the number of VM required less will be the priority as depicted in table 2. Also the load balancer will manage a table of priority of datacenters. The priority will be assigned to datacenters on the basis of VMC. More the VMC a datacenter had more is the priority. This is because higher VMC will lead in increase in overall efficiency of that datacenter and thus more priority.

Virtual Machine count(VMC)	
Total VMs required	Priority
1-3	40%
4-6	30%
7-10	20%
>10	10%

Table 2: Priority assignment based on VMs

Datacenter	
Total VMs	Priority
1-3	10%
4-6	20%
7-10	30%
>10	40%

Table 3: Priority assignment to Datacenter.

5.4 Priority calculation: Here we have considered 4 users U1, U2, U3 and U4 whose parameters are stated in table 3. The

load balancer will calculate the priority of each user task by using table 1 and 2.

Priority is calculated as follows. Let us consider U1. VMC required here is 6 and by table 2 we can see that the priority for VMC=6 is 30%. Again CBQR is 1000001 and by looking in table 1 we can see that when  $CBQR > 1000000$  priority is 60%. Hence total priority to U1 is 90%. Similarly, priorities of all users can be calculated as demonstrated in table 3.

Priority table of user tasks				
UID	VMC	CBQR	Priority	GL
U1	6	1000001	90%	Asia
U2	8	60000	60%	Africa
U3	4	40000	30%	America
U4	7	20000	40%	Asia

Table 4: Priority table of users

Data Center in Asia Region			
DC ID	VMC	VMC available	Priority
DC 1	10	10	40%
DC2	8	8	30%
Dc3	7	7	30%

Table 5: Data center in Asia Region

Data Center in Africa Region			
DC ID	VMC	VMC available	Priority
DC4	9	9	30%
DC5	7	7	30%

Table 6: Data center in Africa Region

Data centre in America Region
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DC ID	VMC	VMC available	Priority
DC 6	4	4	20%
DC7	5	5	20%
Dc8	2	2	10%

Table 7: Data center in America Region

5.5 Datacenter allocation process: The LB now has table of user tasks & datacenter table along with their priority and GL, VMC, priority. The LB will look for the maximum priority task [6][7]. Next it will check mapping. It checks the GL of the task and searches for the cluster of datacenter of that location or nearby location in order to minimize response time. Load balancer will check datacentres according to range of priorities of table 2. Once it finds a datacenter with VMs which satisfies the task need, it will check number of VMs available at that time. If count of VMs present are equal to VMs needed, task will be allocated to that VM otherwise another datacenter is searched. In our example the LB will first process task of U1. It will check its GL which is Asia. Then it will search for the cluster of datacenter that have GL= Asia or nearby Asia. Now the LB will check datacentres according to their range of priorities & will apply best fit to the task [9][10].

#### 5.6 Waiting list

1. Task left idle will be placed in the waiting queue according to their GL and decreasing priorities [5].
2. This waiting list is handled by the load balancer.
3. When execution gets over, it deallocates the VMs and the status is updated in the DC table of that particular GL.
4. Request with higher priority in the waiting list is checked & executed else next task in the waiting list is considered.

#### Proposed Algorithm

Phase 1:

In first phase all VMs in the datacenter are available.

Step 1: Let there be N set of users {U1, U2, Un}.

Step 2: The load balancer will maintain 3 priority assignment tables that is table 1, 2 and 3.

Step 3: Load balancer receives task.

Step 4: Load balancer will calculate the priority of the user tasks using table 1 and 2.

Step 5: LB will maintain a table of datacenters by making cluster on the basis of their GL.

Step 6: The LB will be possessing information about the VMCs in the datacenter and the geographical location of that datacenter.

Step 7: Load balancer will pick up the task with high priority and will check its GL.

Step 8: Load balancer will look for cluster of datacenter with that particular GL.  
 Step 9: After finding cluster of datacenters with that particular GL load will be assigned to datacenter with the minimum number of VMs by applying best fit.  
 Step 10: Update the status of virtual machines available in the datacenters table (table 7, 6, 5).  
 Step 11: Begin the execution.  
 Step 12: Deallocates the VM once the execution is completed and update the status in datacenter table.

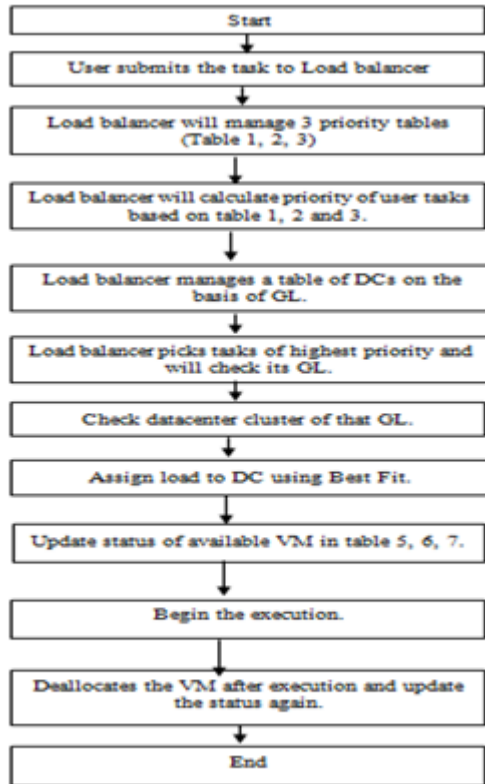


Figure 1 Flowchart of phase 1

Phase 2:  
 VMs in the datacenter are not available.

Step 1: Let there be N set of users {U1, U2... Un}.  
 Step 2: The load balancer will maintain 3 priority assignment tables i.e. table 1, 2 and 3.  
 Step 3: Load balancer receives task.  
 Step 4: Load balancer will calculate the priority of the user tasks using table 1 and 2.  
 Step 5: LB maintains table of datacenters by making cluster depending on GL and will update VMC.  
 Step 6: The load balancer will be possessing information about the Virtual Machine Counts in the datacenter, available virtual machines and the geographical location of that datacenter.

Step 7: Load balancer will pick up the task with high priority and will check its GL.  
 Step 8: Load balancer will look for cluster of datacenter with that particular GL.  
 Step 9: After finding cluster of datacenters with that particular GL, load balancer will look for datacenter with the nearest Virtual Machine Count required. and will also check the available Virtual Machine Counts with datacenter.  
 Step 10: If it satisfies the need of that task, it is allocated to the datacenter. Else repeat step 9 to 10.  
 Step 11: Update the status of virtual machines available in the datacenters table (table 7, 6, 5).  
 Step 12: Begin the execution.  
 Step 13: Deallocates the VM once the execution is completed and again update the status in datacenter table.  
 Step 14: If there is any task which is not allocated then it will be put in the waiting queue according to the decreasing order of priority.  
 Step 15: Once virtual machines are available tasks are picked up from waiting queue according to priority.

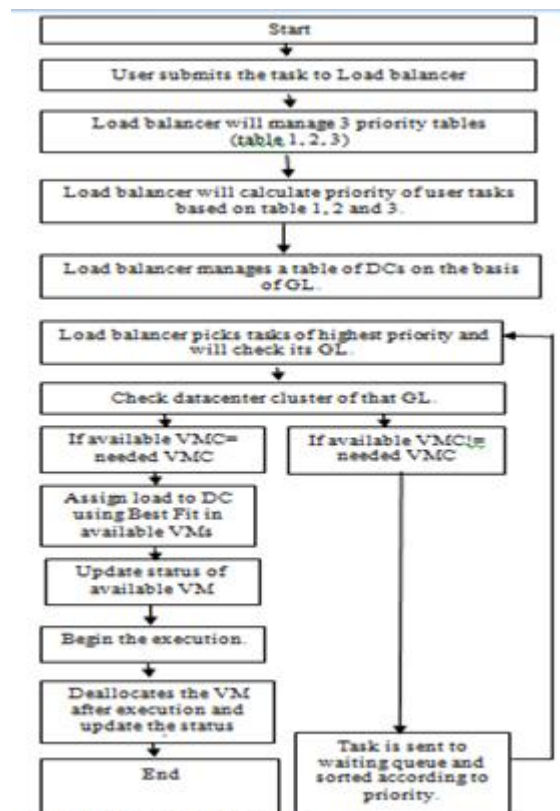


Figure 2 Flowchart for phase 2

#### IV. RESULTS AND DISCUSSION

We have taken cloud analyst tool to represent our results. For simulation three user base (UB) were taken with different priorities. First we put the data center (DC) in the same geographical region of the user base and compute average response time. Next we put data center in different geographical location and compute again the average response time. We find out that the response time for the UB served from DC in same geographical location is far less as compared to UB served from DC in different locations. As per the results we try to minimize the response time [14]. Average response time is reduced if the cloud users will be served from same geographical region and of high priority. UB1 get cloud access first then UB2 and UB3 as shown in table 8. Figure 9 variation in response time based on values of table 8.

S.No	User Base(UB) with priority	Average response time(ms)	
		UB served from Data Centre in same Geographical location	UB served from Data Centre in different Geographical location
1.	UB1 (p1)	50.41	250.98
2.	UB2 (p2)	53.04	299.55
3.	UB3 (p3)	58.29	310.95

Table 8: Results comparison

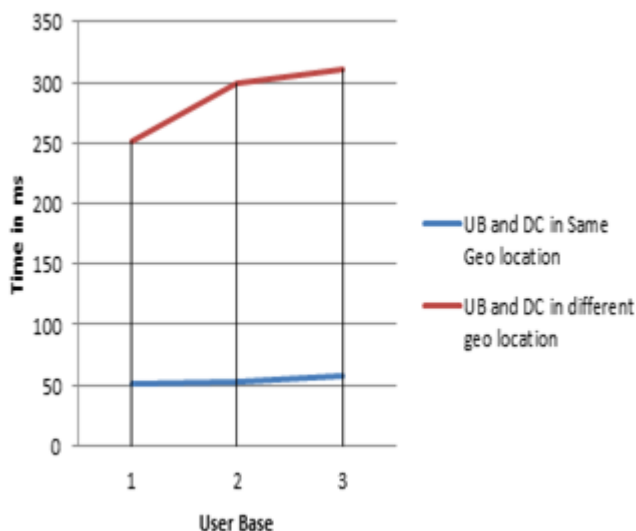


Figure 3: Graph plotted from table 8

#### V. CONCLUSION AND FUTURE SCOPE

In today's situation it is obligatory for cloud service providers to have great relationship with their valued client. To maintain this relationship they need to serve their client needs in most efficient manners.

This paper briefs about the priority based load balancing approach that have a numerous advantages over traditional approach. The whole paper is divided into three basic sections. Section one depicts about the basic of load balancing and different load balancing approach. Section two focuses on proposed methodology and flowchart. Methodology depicts how priority is assigned to the users and datacenter and how these valued customers are served effectively and efficiently. Section three sheds lights on the proposed algorithm steps. When a customer sends a request for a cloud access then that customer can be provided services from that geographical area on the basis of his priority. Serving the customer on the basis of geographical locations in turn reduces the response time and hence increases the efficiency of the cloud.

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